Operating Systems 2 Networking

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Outline

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Introduction

Unix is one of the first operating systems that offered an implementation of network communication. Nowadays, most of the Internet servers run on Linux, the Unix-like operating system. In this lecture a short overview of the Linux kernel network subsystem is given. This subject is complex, so only the most important concepts are presented. The content is split into three parts:

- kernel-level packet processing,
- network device drivers,
- the netfilter implementation.

The part of kernel that is responsible for handling the incoming and outgoing network packets is called the TCP/IP stack. Figure 1 presents the flow of packet processing inside the Linux kernel¹. The Linux kernel network subsystem consists of three parts that corresponds to three layers of ISO/OSI model — the data link layer, the network layer and the transport layer. To send data through a network a user process invokes an appropriate system call that activates the write() method of a file object associated with the process network socket. Depending on the transport protocol used by this socket the write() method calls either the tcp_sendmsg() or udp sendmsg() kernel function. These functions are responsible for building a header of the required protocol.

¹This part of the lecture is based on: https://www.linuxfoundation.org/collaborate/workgroups/networking/network_overview and William Stallings "Operating Systems: Internals and Design Principles", Pearson Education, Inc, London, 2005

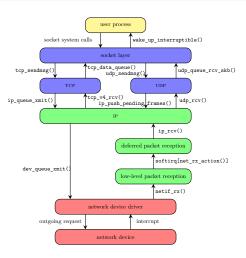


Figure 1: Incoming and outgoing packets processing inside the kernel

The transport protocol header is attached to the data from the user process. Then a function responsible for creating and adding to the packet an IP header is invoked. In case of a UDP packet it is the ip push pending frames() function. For the TCP packet the ip_queue_xmit() is called. The packet with all required headers is passed to a network device with the help of the dev_queue_xmit() function. However, before the packet will be send, its route to the destination host (a computer or other device connected to network) must be established. This is the responsibility of ip_route_output() function that checks caches or (if necessary) routing tables to determine the packet destination. Should the package be sent to other hosts in the network, then it is further processed by the ip output() function.

When a network device receives an incoming packet it usually triggers an interrupt. Optionally two other interrupts may be raised when a transmission of a packet is finished or when a transmission exception occurs.

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There are cases when the network device doesn't trigger the interrupt after reception of a packet. This is explained in the next section where the NAPI is described. The network device driver allocates memory for a packet buffer and sets its pointer to point to the packet IP header. The packet inside the buffer is passed to the netif rx() function which adds it to a queue. All packets from the queue are processed by the ip rcv() function that calls the ip_local_deliver() function. The latter invokes the tcp_v4_rcv() functions for TCP packets or the udp_rcv() function for UDP packets. Next, the functions informing the user process that a packet has arrived are called. In case of TCP packet it is the tcp_data_queue() function and for the UDP protocol it is the udp_queue_rcv() function.

The main data structure used by the kernel network subsystem is the packet buffer called sk buff. The data type of this buffer is struct sk buff. The structure stores not only the received or sent data but also metadata required for processing the packet. The metadata are located in the packet header. The packet buffer is designed to be efficiently transferred from one queue to another. If it is copied then only its header is duplicated. The header has three fields (members) that point to the private headers storing metadata associated with the three layers of ISO/OSI model. The transport_header points to the transport layer header. The network layer header is pointed by the network_header. Finally the mac_layer points to the data link layer header. All packet buffers are stored in a queue implemented as a doubly linked list.

Network Device Drivers

The main data structure used by a network device is a structure of the struct net_device type ². The most important fields of this structure are: the mtu — it specifies the maximum size of a packet that can be transmitted by the device, the flags — it specifies the state of the device, the dev_addr — stores the MAC address, the promiscuity — it is a counter that stores the number of request to set the device in promiscuous mode and the ip_ptr — it points to a variable that stores IPv4 specific data.

Earlier implementations of network device drivers required the device to acknowledge every packet reception by triggering an interrupt. In effect a heavy network traffic could cause a kernel overload. In the 2.6 kernel series a new API for network device drivers was created and named NAPI.

²This part of the lecture is based on http://www.linuxfoundation.org/collaborate/workgroups/networking/network_overview and http://www.linuxfoundation.org/collaborate/workgroups/networking/napi

Network Device Drivers

The New API enables the driver to switch the device into a polling mode and allow it to accumulate a number of incoming packets that are later processed by the kernel. This reduces the number of interrupts triggered by the device and as a consequence lower the kernel load. The driver can also drop some of the packets before they are passed for processing to other parts of the kernel. It is called *packet throttling*. The NAPI requires a buffer in the RAM for DMA transmissions or a hardware support in a form of a DMA *ring*.

The netfilter (the name is an abbreviation of the expression "Network Filter") is a set of function pointers, called *hooks*, that are located in strategic places inside the TCP/IP stack. These pointers can be used for creating firewalls or NAT (Network Address Translation) subsystems. Functions pointed by hooks are usually implemented inside a kernel module³. There are five hooks in the kernel network system:

NF_IP_PRE_ROUTING functions associated with this hook are called when a packet is received,

NF_IP_LOCAL_IN functions associated with this hook perform processing of packets delivered to the host,

NF_IP_FORWARD functions associated with this hook perform processing of packets that should be forwarded to other hosts,

 $^{^3 \}rm http://www.paulkiddie.com/2009/11/creating-a-net filter-kernel-module-which-filters-udp-packets$

```
NF_IP_POST_ROUTING functions associated with this hook perform processing of packages with established route that are intended to be sent,
```

NF_IP_LOCAL_OUT functions associated with this hook perform processing of packets that were sent locally.

Each function associated with any of the hooks can perform any operation on a packet that is necessary, but it has to eventually return one of the following values:

```
NF_ACCEPT the packet is accepted for further processing,
NF_DROP the packet is rejected,
```

NF_REPEAT the function call should be repeated for this package,

NF_STOLEN the function "steals" the packet, which means that this packet will be processed in a different way than the other packages,

NF_QUEUE the packet is added to a queue from where it will be transfered to the user-space,

NF_STOP processing of the packet is stopped.

A single function associated with a hook is represented by a structure of the nf_hook_ops type. The definition of the type is given in the Listing 1. The list field allows these structures to be stored in a linked list. The hook field is a pointer to the packet processing function. The dev field is a pointer to a structure that represents the network device. The priv member is a pointer to an area of the memory that stores private data of the packet processing function. The pf field stores the identifier of protocol family. Packets of that protocols will be processed by the function. The hooknum field stores the hook number and the priority field stores the hook priority that specifies the order in which the hooks are performed (the NF_IP_PRI_FIRST constant defines the highest priority).

```
struct nf hook ops
   {
            struct list_head
                                       list;
3
            nf_hookfn
                                       *hook;
            struct net_device
                                       *dev;
            void
                                       *priv;
            u_int8_t
                                       pf;
            unsigned int
                                       hooknum;
            int
                                       priority;
9
   };
10
```

Listing 1: The definition of the struct nf_hook_ops

Structures of the struct nf_hook_ops type are registered with the use of the nf_register_hook() function and unregistered with the help of the nf_unregister_hook() function. Each packet processing function has to return a number of the unsigned int type and take three arguments: an address of its private data (it is passed by the void * type parameter), the address of a packet buffer (the buffer is of the struct sk_buff type) and finally an address of a structure that stores the state of the hook. This structure is of the struct nf_hook_state type.

Questions

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THE END

Thank You for Your attention!